МАТЕМАТИЧНЕ ТА КОМП'ЮТЕРНЕ МОДЕЛЮВАННЯ

MATHEMATICAL AND COMPUTER MODELING

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ANALYSIS OF THE RESULTS OF SIMULATION MODELING OF THE INFORMATION SECURITY SYSTEM AGAINST UNAUTHORIZED ACCESS IN SERVICE NETWORKS

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ABSTRACT

Context. An analysis of the service network shows that insufficient information security in service networks is the cause of huge losses incurred by corporations. Despite the appearance of a number of works and materials on standardization, there is currently no unified system for assessing information security in the field of information security. It should be noted that existing methods, as well as accumulated experience in this area, do not completely overcome these difficulties. This circumstance confirms that this problem has not yet been sufficiently studied and, therefore, remains relevant. The presented work is one of the steps towards creating a unified system for assessing information security in service networks.

Objective. Development of an algorithm and simulation model, analysis of simulation results to determine the key characteristics of the Information Security System, providing the capability for complete closure, through the security system, of all potential threat channels by ensuring control over the passage of all unauthorized access requests through defense mechanisms.

Method. To solve the problem, a simulation method was applied using the principles of queuing system modeling. This method makes it possible to obtain the main characteristics of the Information Security System from the unauthorized access with a limited amount of buffer memory.

Results. Algorithms, models, and methodology have been developed for the development of Information Security System from unauthorized access, considered as a single-phase multi-channel queuing system with a limited volume of buffer memory. The process of obtaining model results was implemented in the General Purpose Simulation System World modelling system, and comparative assessments of the main characteristics of the Information Security System were carried out for various laws of distribution of output parameters, i.e., in this case, unauthorized access requests are the simplest flows, and the service time obeys exponential, constant, and Erlang distribution laws.

Conclusions. The conducted experiments based on the algorithm and model confirmed the expected results when analyzing the characteristics of the Information Security System from the unauthorized access as a single-phase multi-channel queuing system with a limited waiting time for requests in the queue. These results can be used for practical construction of new or modification of existing Information Security System s in service networks of objects of various purposes. This work is one of the approaches to generalizing the problems under consideration for systems with a limited volume of buffer memory. Prospects for further research include research and development of the principles of hardware and software implementation of Information Security System in service networks.

KEYWORDS: unauthorized access, information security systems, information security, queuing systems, defense mechanism, simulation modeling.

ABBREVIATIONS

BM is a Buffer Memory;

DM is a Defense Mechanism; GPSS World is a General Purpose Simulation System (latest version of GPSS); ISS is an Information Security System; QS is a Queuing System; UA is an Unauthorized Access.

NOMENCLATURE

AVE.C is an Average Queue Length; CUM.% is a Cumulative Percentage, expressed as a percentage of the total number of random values; © Ismailov B. G., 2024 DOI 10.15588/1607-3274-2024-2-4 ENTRIES is a number of requests in DM;

FREQUENCY is the number of random values falling within the given interval;

 L_a is an average queue length;

 L^0 is the permissible limit values L_a ;

M is a mathematical expectation symbol;

MEAN is a mean value of the corresponding random variable;

N is a number of DMs in ISS;

 N_0 is the permissible limit values N;



 p_1 is a probability of blocking UA requests;

 p_2 is a probability of UA requests bypassing protected resources;

RANGE is a lower and upper bound of the frequency interval;

RETRY is a number of requests waiting for the fulfillment of a specific condition depending on the state of this table;

STD.DEV is a Standard Deviation of the random variable:

 T_U is a time of requests' stay in the system;

 T_W is a time of requests' waiting in the queue;

UTIL is a Utilization Coefficient of DM;

 λ is an intensity of various threats at the entrance of ISS:

 λ_0 is the permissible limit values λ ;

 μ is an intensity of servicing UA requests;

 μ_0 is the permissible limit values μ ;

 τ_0 is a service delays;

 ρ is a normalized intensity.

INTRODUCTION

This work is dedicated to approaches in researching ISS in service networks, addressing security issues characteristic of systems with limited BM capacities. When addressing security issues in service networks, the primary determinant is the security class of the network, defining a set of DMs that constitute the hardware or software part implemented in the network. In service networks, intentional UA requests are often received alongside regular requests, targeting confidential information from illegal users, which can lead to network disruptions. It should be noted that DMs, influencing the entire information security process, may operate in constant information interaction with other elements of the ISS. The operation of DMs is described by four possible states: operational, non-operational, diagnosed, and restored. In ISS, the possibility of an undesirable event related to the reliability characteristics of DMs, leading to various types of losses, is considered a risk. However, approaches associated with the risk arising from the reliability characteristics of DMs are not considered in this work, i.e., it is assumed that all DMs are reliable.

The task related to the security problem in service networks is addressed by examining the ISS, ensuring the complete closure of all possible channels of threat manifestation through the security system. This is achieved by controlling the passage of all UA requests through DM.

The object of the study is an ISS against UA with a limited amount of BM in service networks.

The subject of the study is to determine the structure of the object, i.e. determination of the main characteristics of the system - security of information from UA with a limited amount of BM in service networks.

The goal of the work is to develop an algorithm and simulation model, analyze the results of the simulation

© Ismailov B. G., 2024 DOI 10.15588/1607-3274-2024-2-4 model, allowing us to determine the main characteristics of an ISS against UA with a limited amount of BM in service networks.

1 PROBLEM STATEMENT

The structure of ISS with limited BM is considered (Fig. 1), where all input streams are directed to DM for servicing. As noted earlier, the security system allows for the complete closure of all possible channels of threat manifestation by controlling the passage of all UA requests through DM. It is assumed that the examined ISS structure ensures maximum information security for service networks. This structure constitutes a hardware and software complex interacting with random event streams, which are conditioned by the actions of attackers, improper access rights distribution, unauthorized software usage, as well as errors in identification and authentication software and technical complexes.



The assumption is that the intruder (attacker, UA requests) at the system's entrance generates various threats with intensity λ . The ISS consists of N DMs that introduce delays $\tau_0 = 1/\mu$ in service. If we consider the intruder block as an information source and DMs as devices operating in parallel, the mathematical model of the ISS can be regarded as a single-phase, multi-channel QS with limited BM. Taking into account the complex nature of UA request servicing (filtering UA requests, detection and classification of UA attempts, blocking or allowing UA requests to access protected resources, etc.), Poisson formulas are suggested as the probability loss function for UA requests due to system overload in [1]:

$$P(\lambda,\mu,N) = \sum_{j=N}^{\infty} \left(\rho^j / j! \right) e^{-\rho}.$$

Then, the problem of determining the optimal values of ISS characteristics can be formulated as the minimization of the mathematical expectation of the probability loss function for UA requests due to system overload:



$$M\left[\sum_{j=N}^{\infty} \left(\rho^j / j!\right) e^{-\rho}\right] \to \min$$

with $\lambda \ge \lambda_0$, $\mu \ge \mu_0$, $N \ge N_0$, $L_q \le L^0$,

where $\rho = \lambda/\mu$. Problems related to insufficient information security in service networks, and the task of determining optimal ISS characteristics against UA for various cases, have been considered and analytically solved in [1] and optimal values for QS characteristics with and without waiting requests in the queue have been obtained.

However, for a detailed analysis of ISS characteristics against UA across a wide range of input and output parameters, it is preferable to utilize simulation modeling methods, considering it as a single-phase, multi-channel QS both with and without waiting. Given the volume of obtained results from the simulation model, we will limit the discussion here to the analysis of the simulation model results for a QS with limited waiting requests, encompassing a broad range of input and output parameters.

Thus, based on the presented structure of the ISS, the task in this work is to analyze the results of simulation modeling of a single-phase multi-channel QS with limited BM. To achieve this, using simulation modeling, it is necessary to determine the structural and temporal characteristics of the ISS within the specified values of concurrently operating service devices (DMs).

2 LITERATURE REVIEW

Analysis and accumulated experience demonstrate that insufficient information security in service networks leads to significant losses for corporations. This underscores the high importance of the information security problem. An analysis of the current state of the issue in the field of information security and the development of ISS reveals serious challenges, largely stemming from the absence of a unified system for assessing information security. Such a system would enable a quantitative evaluation during the design and operation of service networks [2–8]. It is worth noting, that due to a lack of sufficient experience in designing ISS, tasks related to its construction must be addressed at the early stages of service network design.

Currently, given the increasing number of scientific studies and companies specializing in information security in service networks, this problem is insufficiently explored [4–11, 13]. It should be noted that one of the most obvious causes of ISS violation is intentional UA requests for confidential information by illegal users, followed by undesirable manipulations with this information [1, 2,12]. The effectiveness of information security protection in service networks is primarily determined by the service network's security class [1, 2, 11, 14–16], which defines the set of DMs implemented in the network.

In [1], due to the fact that the security system fails to completely close all possible threat channels, a structure for the ISS was proposed. Unlike existing structures, in this framework, each input stream is provided with a DM for maintenance.

In the work [1], a structure for the ISS with losses is proposed, featuring both limited and unlimited BM. This structure ensures maximum information security in service networks by controlling the passage of all UA requests through DMs. In contrast to [1], an analysis of the simulation model results for ISS with limited BM is conducted here, encompassing a broad range of input and output parameters.

3 MATERIALS AND METHODS

To determine the characteristics of the ISS that allow it to operate within limited resources, it is assumed that the input flow of information, i.e., UA requests, is Poisson distributed, and the service time follows exponential, constant, and Erlang distribution laws. Algorithms for the simulation model of the service process have been developed for three cases to adequately describe the operation of the ISS against UA:

1. Incoming requests to the ISS and service time follow an exponential distribution.

2. Incoming requests to the ISS follow an exponential distribution, while the service time follows a uniform distribution.

3. Incoming requests to the ISS follow an exponential distribution, while the service time follows an Erlang distribution.

The developed algorithm for the operation of the ISS against UA includes the following steps:

- setting the minimum permissible limit values for the number of concurrently operating service devices (DMs) and the maximum permissible limit values for the queue length, defining the BM volume;

- to conduct a detailed analysis of the properties of the investigated system, a table structure is organized for queue waiting time and request residence time. An upper limit for the first frequency interval is specified, along with the values for all other frequency intervals and the quantity of frequency intervals. The goal here is to build histograms of the probability density function for the waiting time in the queue and the residence time of requests in the system based on the accumulation of the frequency of occurrence of random variables within the specified frequency intervals;

- when a request is received, the system checks for available places in the queue. If there is no available space in the queue, the request is rejected and exits the system;

- otherwise, if all DMs are occupied, the UA request waits in the system's BM queue until one of the DMs becomes available provided there is free space in the BM.

– upon the release of one of the DMs, the UA request enters this available DM, and the process of filtering UA requests, detecting, and classifying UA attempts takes place. As a result, the initial UA stream is thinned out with certain probabilities p_1 , $p_2 = 1 - p_1$ forming an output stream, in other words, with a probability of p_1 block-

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ing occurs, while with a probability of p_2 UA requests are allowed to pass through to the protected resources.

Note 1. The values of probabilities p_1 , p_2 are determined based on statistical analysis.

Based on the proposed algorithm covering three cases of ISS functioning against UA as single-phase, multichannel QS with a limited buffer size, simulation models were developed using the GPSS World modeling language. For $N = \overline{2,5}$ during the simulation the model allows you to determine such characteristics as ENTRIES, UTIL, AVE.C, MEAN, STD.DEV, RANGE, RETRY, FREQUENCY, CUM.%.

4 EXPERIMENTS

Based on the execution of the simulation model for the average values of real data, with $N = \overline{2,5}$, $\lambda = 1/3500 \text{ ms}$ and $\mu = 1/1700 \text{ ms}$ results were obtained for three cases:

1. Incoming requests to the ISS and service time follow an exponential distribution.

In the first case, the results of a simulation model of the functioning of the information system were obtained – reports and histograms of the distribution densities of the residence time $T_{\rm U}$ and waiting time $T_{\rm W}$ of requests, with

 $N = \overline{2,5}$ (see Appendix A, Fig. A.1–A.8).

Based on the obtained reports, Table 1 was created, providing the dynamics of changes in the number of requests in the DM, the average queue length, and the utilization coefficient of the DM depending on the number of DM (N) during the modeling period for the first case.

Table 1 – Dynamics of changes in characteristics depending on the number of DMs for the first case

The	The number of	The aver-	The utilization
number of	requests in the	age queue	coefficient of
DM	DM	length	the DM
2	90266	1.866	0.933
3	99605	2.047	0.682
4	99989	2.056	0.514
5	100002	2.070	0.414

The analysis of the dynamics of these parameters shows that with an increase in the number of DM from 2 to 5:

- the number of requests in DM increases, with a difference of 9736 requests;

- the average queue length increases, with a difference of 0.204;

- the utilization coefficient of DM decreases, with a difference of 0.519.

In the models, 10 frequency intervals were chosen for building histograms, and the length of frequency intervals was selected as 0.0004 time units for waiting time in the queue and 0.0008 time units for the service time. The analysis conducted shows that in the first case, with a change in the number of DM from 3 to 5, the characteristics of the density distribution of the residence time

© Ismailov B. G., 2024 DOI 10.15588/1607-3274-2024-2-4 T_U and waiting time T_W of requests do not change.

Note 2. For clarity of histograms, it is desirable to have a large number of frequency intervals. To obtain an objective picture, it is necessary to have a large sample of random variables, which is not always possible and feasible.

Note 3. The values of interval lengths and the number of frequency intervals are selected experimentally during several runs of the simulation model or based on assumed values of the mean and standard deviation of the corresponding random variable.

2. The requests entering the ISS follow an exponential distribution, while the service time adheres to a uniform distribution.

In the second case, the results of a simulation model of the functioning of the ISS were obtained – reports and histograms of the distribution densities of the residence time and waiting time of requests, with $N = \overline{2,5}$ (see Appendix B, Fig. B.1–B.8).

Based on the obtained reports, Table 2 was created, providing the dynamics of changes in the number of requests in the DM, the average queue length, and the utilization coefficient of the DM depending on the number of DM (N) during the modeling period for the second case.

Table 2 – Dynamics of changes in characteristics depending on the number of DMs for the second case

The num-	The number of	The aver-	The utilization
ber of DM	requests in the	age queue	coefficient of
	DM	length	the DM
2	93922	1.932	0.966
3	99980	2.061	0.687
4	100002	2.056	0.514
5	100002	2.053	0.411

The analysis of the dynamics of these parameters shows that with an increase in the number of DM from 2 to 5:

- the number of requests in DM increases, with a difference of 6080 requests;

- the average queue length increases, with a difference of 0.121;

- the utilization coefficient of DM decreases, with a difference of 0.555.

The analysis conducted shows that in the first case, with a change in the number of DM from 3 to 5, the characteristics of the density distribution of residence time T_U and waiting time T_W of requests do not change.

3. The incoming requests to the ISS follow an exponential distribution, while the service time follows an Erlang distribution.

In the third case, the results of the ISS simulation model were obtained – reports and histograms of the density distribution of the residence time $T_{\rm U}$ and the waiting

time T_W of requests at $N = \overline{2,5}$ (see Appendix C, Fig. C.1–C.8).



Based on the obtained reports, Table 3 was created, providing the dynamics of changes in the number of requests in the DM, the average queue length, and the utilization coefficient of the DM depending on the number of DM (N) during the modeling period for the third case.

Table 3 – Dynamics of changes in characteristics depending on the number of DMs for the third case

i	the number of Divis for the third ease										
The num-	The number	The aver-	The utilization								
ber of DM	of requests in	age queue	coefficient of the								
	the DM	length	DM								
2	91907	1.898	0.949								
3	99633	2.053	0.684								
4	100003	2.061	0.515								
5	100002	2.058	0.412								

The analysis of the dynamics of these parameters shows that with an increase in the number of DM from 2 to 5:

- the number of requests in DM increases, with a difference of 8095 requests;

- the average queue length increases, with a difference of 0.159;

- the utilization coefficient of DM decreases, with a difference of 0.537.

The analysis conducted shows that in the first case, with a change in the number of DM from 3 to 5, the characteristics of the density distribution of the residence time T_U and waiting time T_W of requests do not change.

Based on Tables 1–3, the dynamics of changes in the differences in the number of requests in the DM, average queue length, and the utilization coefficient of the DM for three cases with $N = \overline{2,5}$, and the nature of these differences are presented in Fig. 2–4.





Figure 2 – The nature of the change in the differences in the number of requests in the DM for three cases with $N = \overline{2,5}$

AVE.C



age queue length for three cases with $N = \overline{2,5}$



Figure 4 – The nature of changes in the differences in the coeffi-

cients of use of DM for three cases with $N = \overline{2,5}$

The results obtained from Table 1–3 and Fig. 2–4 show that with an increase in the number of DM from 2 to 5 in three cases:

- the nature of the change in the differences in the number of requests in the DM is 9736, 6080 and 8095;

- the nature of the change in the differences in the average queue length is 0,204; 0,121 and 0.163;

- the nature of the change in the differences in the utilization coefficient DM is 0.519; 0.555 and 0.537.

CONCLUSIONS

The current task of developing an algorithm and simulation model, along with the analysis of simulation model results to determine the key characteristics of the ISS, is being addressed. This aims to provide the capability for complete closure, through the security system, of all potential threat channels by ensuring control over the transition of all UA through the DM.

The scientific novelty of the obtained results lies in the fact that, for the first time, algorithms, and simulation models, as well as a methodology for developing the ISS, have been proposed and developed based on the analysis of structural and temporal characteristics of ISS from UA. This includes treating it as a single-phase multi-channel queueing system with limited waiting time in the queue across a wide range of input and output parameters. The experiments conducted using the algorithm and model

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confirmed the expected results when analyzing the characteristics of the ISS from UA.

The practical significance of the results lies in their applicability for the practical construction of new or modification of existing ISS in networks for various purposes. This work represents one of the approaches to generalizing the considered problems for systems with a limited BM.

Prospects for further **research** include the exploration and development of hardware and software implementation principles for ISS from UA with a limited BM in service networks.

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Appendix A

The first case's simulation model's results of the ISS's operation.

QUEUE CH_1	MAX CONT. 10 1	ENTRY E 90266	NTRY (0) 10069	AVE.CON 4.82	NT. AVE.TIN 6 0.00	ME AVE.(-0 02 0.00) RETRY 2 O
STORAGE UZEL	CAP. REM. 2 0	MIN. MA O	X. EN 2 90	TRIES AVI	L. AVE.C. 1.866	UTIL. RETRY 0.933 0	DELAY 0
TABLE T W	MEAN ST 0.002 0	D.DEV.	1	RANGE	RET	TRY FREQUENC	Y CUM.&
-					0.000	19250	21.33
		0	.000 .		0.001	9868	32.26
		0	.001 .		0.001	10215	43.58
		0	.001 .		0.002	10616	55.34
		0	.002 .		0.002	10424	66.88
		0	.002 .	-	0.002	9052	76.91
		0	.002 .		0.003	7398	85.11
		0	.003 .	-	0.003	5324	91.01
		0	.003 .		0.004	3461	94.84
		0	.004 .			4657	100.00
T_U	0.002 0	.001			0)	
					0.001	14294	17.58
		0	.001 .		0.002	17199	38.73
		0	.002 .		0.002	18110	61.00
		0	.002 .	-	0.003	15313	79.83
		0	.003 .		0.004	9382	91.37
		0	.004 .	-	0.005	4420	96.81
		0	.005 .		0.006	1704	98.90
		0	.006 .	-	0.006	629	99.68
		0	.006 .		0.007	206	99.93
		0	.007 .			57	100.00





Figure A.2 – Histograms of the probability density functions of the residence time T_U and waiting time T_W of requests N=2

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QUEUE CH_1	MAX CONT. ENTRy 10 1 99605	ENTRY(0) AVE 53200 0	.CONT. AVE.TIME .951 0.000	AVE.(-0) REIRY 0.001 0
STORAGE UZEL	CAP. REM. MIN. 3 0 0	MAX. ENTRIES 3 99605	AVL. AVE.C. UT 1 2.047 0.	IL. RETRY DELAY 682 0 0
TABLE	MEAN STD.DEV	. RANGE	RETRY	FREQUENCY CUM. &
TW	0.000 0.000		0	
-		-	0.000	75800 76.10
		0.000 -	0.001	11400 87.55
		0.001 -	0.001	6100 93.67
		0.001 -	0.002	3356 97.04
		0.002 -	0.002	1679 98.73
		0.002 -	0.002	809 99.54
		0.002 -	0.003	285 99.82
		0.003 -	0.003	109 99.93
		0.003 -	0.004	46 99.98
		0.004 -		20 100.00
T_U	0.001 0.001		0	
			0.001	51706 57.64
		0.001 -	0.002	24423 84.87
		0.002 -	0.002	9341 95.28
		0.002 -	0.003	3045 98.68
		0.003 -	0.004	828 99.60
		0.004 -	0.005	272 99.90
		0.005 -	0.006	68 99.98
		0.006 -	0.006	17 100.00
		0.006 -	0.007	3 100.00
Eiguna A 2	Ero amont of	the remark	for the mode	1 \dots $M = 2$





Figure A.4 – Histograms of the probability density functions of the residence time T_U and waiting time T_W of requests N=3

QUEUE CH_1	MAX C0 10	ONT. O	ENTR) 99989	9 8130	(0) AVE 9 0	.CONT.	AVE.TI	ME AVE.(00 0.	-0) RETRY 000 0
STORAGE UZEL	CAP. 1 4	REM. 3	MIN. 0	MAX. 4	ENTRIES 99989	AVL. 1	AVE.C. 2.056	UTIL. RET 0.514	RY DELAY 0 0
TABLE T W	MEAN 0.000	SI	ID.DE	7.	RANGE		RE	TRY FREQUE	NCY CUM.&
					-	0.	000	94942	94.95
				0.000	- 1	ο.	001	3638	98.59
				0.001	-	0.	001	1052	99.64
				0.001	-	0.	002	258	99,90
				0.002	-	0.	002	80	99.98
				0.002	-	0.	002	16	100.00
				0.002	-	0.	003	3	100.00
TU	0.001	(0.001					0	
					-	0.	001	63515	70.46
				0.001	-	0.	002	19535	92.13
				0.002	-	0.	002	5254	97.96
				0.002	-	0.	003	1343	99.45
				0.003	- 1	0.	004	366	99.85
				0.004	-	0.	005	106	99.97
				0.005	-	0.	006	25	100.00
				0.000	- :	0.	006	1	100.00
				0.006		0.	007	1	100.00
D ' A C D			C /	1		.1	1	1 14	37 4

Figure A.5 – Fragment of the report for the model with N = 4





Figure A.6 - Histograms of the probability density functions of the residence time T_U and waiting time T_W of requests N = 4

QUEUE CH_1	MAX CONT 9 (. ENTRY	ENTRY 93071	(0) AVE. L 0.	CONT. AV 051	E.TIME 0.000	AVE.(-0) 0.000	RETRY 0
STORAGE UZEL	CAP. REN 5 S	4. MIN. 1 8 0	MAX. H 5 1	ENTRIES	AVL. AV 1 2.	E.C. UTI 070 0.4	IL. RETRY	DELAY 0
TABLE	MEAN	STD.DEV		RANGE		RETRY	FREQUENCY	CUM.&
·_"	0.000	0.000		-	0.000		99001	99.00
			0.000	-	0.001		817	99.82
			0.001	-	0.001		167	99.98
			0.001	-	0.002		15	100.00
			0.002	-	0.002		2	100.00
T_U	0.001	0.001				0		
			_	-	0.001		66085	73.36
			0.001	-	0.002		17898	93.22
			0.002	-	0.002		4488	98.20
			0.002	-	0.003		1185	99.52
			0.003	-	0.004		318	99.87
			0.004	-	0.005		94	99.98
			0.005	-	0.006		20	100.00
			0.006	-	0.006		1	100.00

Figure A.7 – Fragment of the report for the model with N = 5



Figure A.8 - Histograms of the probability density functions of the residence time T_U and waiting time T_W of requests N = 5

Appendix **B**

Results of the second case's simulation model of the ISS's operation.

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QUEUE CH_1	MAX CONT 10 10	ENTRY 93922	ENTRY (5263	(0) AVE 5	.CONT. AV	0.002	AVE.(-0 0.00) RETRY 2 O
STORAGE UZEL	CAP. REM. 2 0	. MIN. 1 0	MAX. E 2	NTRIES 93912	AVL. AV 1 1.	/E.C. UT1 .932 0.9	IL. RETRY 966 0	DELAY 10
TABLE	MEAN	STD.DEV		RANGE		RETRY	FREQUENC	Y CUM.&
TW	0.002	0.001				0		
-				-	0.000)	13253	14.11
			0.000	-	0.001	L	10024	24.79
			0.001	-	0.003	1	11128	36.64
			0.001	-	0.002	2	12288	49.72
			0.002	-	0.002	2	13158	63.73
			0.002	-	0.002	2	13948	78.58
			0.002	-	0.003	3	14840	94.39
			0.003	-	0.003	3	5273	100.00
T_U	0.002	0.001				0		
			_	-	0.003	1	8029	9.49
			0.001	-	0.002	2	18103	30.90
			0.002	-	0.003	2	22165	57.11
			0.002	-	0.003	3	25273	86.99
			0.003	-	0.004	•	11004	100.00
Figure B.1 – Fragment of the report for the model with $N = 2$								



Figure B.2 - Histograms of the probability density functions of the residence time T_U and waiting time T_W of requests



0.001 0.002 0.002 0.003 65090 23406 1438 32 0.001 0.001 Figure B.3 – Fragment of the report for the model with N = 3



Figure B.4 - Histograms of the probability density functions of the residence time T_U and waiting time T_W of requests





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QUEUE CH_1	MAX CONT. ENTR 8 0 10000	Y ENTRY(2 82358	0) AVE. 0.	CONT. AVE	.TIME 0.000	AVE.(-0 0.00) RETRY 0 0
STORAGE UZEL	CAP. REM. MIN. 4 2 0	MAX. E 4 1	NTRIES 00002	AVL. AVE 1 2.0	.C. UT	IL. RETRY 514 O	DELAY 0
TABLE T W	MEAN STD.DE	v.	RANGE		RETRY	FREQUENC	Y CUM.♥
			-	0.000		98663	98.66
		0.000	-	0.001		1293	99.95
T_U	0.001 0.000				0		
_		_	-	0.001		84401	93.84
		0.001	-	0.002		5542	100.00
		0.002	-	0.002		3	100.00

Figure B.5 – Fragment of the report for the model with N = 4



Figure B.6 – Histograms of the probability density functions of the residence time T_U and waiting time T_W of requests N = 4

QUEUE CH_1	MAX CONT. ENT 7 0 1000	RY ENTRY(02 93823	0) AVE.CONT 0.029	. AVE.TIME 0.000	AVE.(-0) RETRY 0.000 0	
STORAGE UZEL	CAP. REM. MIN 5 3 0	. MAX. E 5 1	NTRIES AVL. 00002 1	AVE.C. UT 2.053 0.4	IL. RETRY DELAY	
TABLE T_W	MEAN STD.D	EV. 0	RANGE	RETRY 0	FREQUENCY CUM.&	
		0.000	- 0	0.000	99894 99.89 107 100.00	
T_U	0.001 0.00	0.001	- 0	0.001	1 100.00	
		o	- 0	.001	88876 98.71	
	_	0.001	- (.002	1105 100.00	

Figure B.7 – Fragment of the report for the model with N = 5



Figure B.8 – Histograms of the probability density functions of the residence time T_U and waiting time T_W of requests N = 5

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Appendix C

Results of the third case's simulation model of the ISS's operation.

QUEUE CH_1	MAX CONT. ENTRY 10 6 91907	ENTRY(0) AVE. 7758 4.	CONT. AVE.TIME 921 0.002	AVE.(-0) RETRY 0.002 0
STORAGE UZEL	CAP. REM. MIN. 2 0 0	MAX. ENTRIES 2 91902	AVL. AVE.C. UT 1 1.898 0.	IL. RETRY DELAY 949 0 5
TABLE T W	MEAN STD.DEV 0.002 0.001	. RANGE	RETRY	FREQUENCY CUM.&
-		-	0.000	16613 18.08
		0.000 -	0.001	10031 28.99
		0.001 -	0.001	11184 41.16
		0.001 -	0.002	11418 53.59
		0.002 -	0.002	11645 66.26
		0.002 -	0.002	10743 77.95
		0.002 -	0.003	8832 87.56
		0.003 -	0.003	5696 93.76
		0.003 -	0.004	3276 97.32
		0.004 -		2463 100.00
T_U	0.002 0.001	-	0	
			0.001	11476 13.84
		0.001 -	0.002	18300 35.92
		0.002 -	0.002	20131 60.20
		0.002 -	0.003	18139 82.08
		0.003 -	0.004	10430 94.66
		0.004 -	0.005	3405 98.76
		0.005 -	0.006	851 99.79
		0.006 -	0.006	161 99.98
		0.006 -	0.007	13 100.00

Figure C.1 – Fragment of the report for the model with N = 2



Figure C.2 – Histograms of the probability density functions of the residence time T_U and waiting time T_W of requests N=2

QUEUE CH_1	MAX 10	CONT.	ENTR) 99833	ENTRY 5315	(0) AVE. 2 0.	.CONT. AN	0.000	AVE.(-0 0.00) RETRY 0 0
STORAGE UZEL	CAP. 3	REM. 1	MIN. 0	MAX. 3	ENTRIES 99833	AVL. AV 1 2.	7E.C. UI .053 0.	IL. RETRY 684 0	DELAY 0
TABLE T W	MEAN	5 0	TD.DEV	<i>.</i>	RANGE		RETRY	FREQUENC	Y CUM.&
					-	0.000) -	78532	78.66
				0.000	-	0.001		12483	91.17
				0.001	-	0.001		5268	96.44
				0.001	-	0.002	2	2232	98.68
				0.002	-	0.002	2	932	99.61
				0.002	-	0.002	2	284	99.90
				0.002	-	0.003	3	81	99.98
				0.003	-	0.003	3	14	99.99
				0.003	-	0.004	1	7	100.00
T_U	0.00	1	0.001				0		
				_	-	0.001	L.	52246	58.16
				0.001	-	0.002	2	28786	90.21
				0.002	-	0.002	2	7193	98.21
				0.002	-	0.003	3	1388	99.76
				0.003	-	0.004	1	187	99.97
				0.004	-	0.005	5	26	100.00
				0.005	-	0.006	5	4	100.00
Figure C.3 –	Frag	ame	nt of	the r	eport	for the	mode	el with	N = 3





Figure C.4 – Histograms of the probability density functions of the residence time T = U and waiting time T = W of requests N = 3

QUEUE CH_1	MAX CONT 10 3	. ENTRY 100003	ENTRY (81249	0) AVE 0	.CONT. AVE. .166 0	TIME .000	AVE.(-0) 0.000	RETRY 0
STORAGE UZEL	CAP. REM 4 0	. MIN. 1	MAX. E 4 1	NTRIES	AVL. AVE. 1 2.06	C. UTI 1 0.5	L. RETRY	DELAY 2
TABLE T W	MEAN	STD.DEV		RANGE		RETRY	FREQUENCY	CUM
				-	0.000		96173	96.17
			0.000	-	0.001		3194	99.37
			0.001	-	0.001		509	99.88
			0.001	-	0.002		91	99.97
			0.002	-	0.002		26	99,99
			0.002	-	0.002		7	100.00
T_U	0.001	0.000				0		
			_	-	0.001		64240	71.40
			0.001	-	0.002		22423	96.33
			0.002	-	0.002		2962	99.62
			0.002	-	0.003		308	99.96
			0.003	-	0.004		32	100.00
			0.004	-	0.005		2	100.00

Figure C.5 – Fragment of the report for the model with N = 4



Figure C.6 - Histograms of the probability density functions of the residence time T_U and waiting time T_W of requests N = 4

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QUEUE CH_1	MAX CC 10	0 100	TRY EN1 002 93	RY(0) 3 516	VE.CONT. 0.040	AVE.TIM 0.00	E AVE.(- 0 0.0	0) RETRY 00 0			
STORAGE UZEL	CAP. F	3 0	N. MAX. 5	ENTR: 10000	IES AVL. 02 1	AVE.C. 2.058	UTIL. RETR 0.412 0	Y DELAY			
TABLE	MEAN	STD.	DEV.	RAI	NGE	RET	RY FREQUEN	ICY CUM.&			
			0.0	- 500	0. 0.	.000	99406 535	99.40 99.94			
			0.0	001 -	0.	.001	54	99.99			
T_U	0.001	0.0	00								
			0.0	001 -	0.	.001	20334	97.03			
			0.0	02 -	0.	.002	2419	99.71			
			0.0	02 -	0.	.003	239	99.98			
			0.0	03 -	0.	.004	18	100.00			
			0.0	04 -	0.	.005	1	100.00			
Figure C.7 – Fragment of the report for the model with $N = 5$											



Figure C.8 - Histograms of the probability density functions of the residence time T_U and waiting time T_W of requests N = 5

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АНАЛІЗ РЕЗУЛЬТАТІВ ІМІТАЦІЙНОГО МОДЕЛЮВАННЯ СИСТЕМИ БЕЗПЕКИ ІНФОРМАЦІЇ ВІД НЕСАНКЦІОНОВАНОГО ДОСТУПУ У МЕРЕЖАХ ОБСЛУГОВУВАННЯ

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АНОТАЦІЯ

Актуальність. Аналіз мережі обслуговування показує, що недостатня захищеність інформації в мережах обслуговування є причиною великих збитків, завданих корпораціями. Незважаючи на появу низки робіт, та матеріалів зі стандартизації, в даний час єдина система оцінки захищеності інформації в галузі інформаційної безпеки відсутня. Слід зазначити, що існуючі методи, і навіть накопичений досвід у цій галузі неможливо повністю подолати ці труднощі. Ця обставина підтверджує, що дана проблема ще недостатньо вивчена і, отже, зберігає свою актуальність. Представлена робота є одним із кроків на шляху створення єдиної системи оцінки безпеки інформації у мережах обслуговування.

Мета роботи. Розробка алгоритму та моделі імітації, аналіз результатів моделі імітації для визначення основних характеристик системи безпеки інформації, що надають можливість повного закриття, за допомогою системи безпеки, всіх можливих каналів прояву загроз, шляхом забезпечення контролю переходу всіх запитів несанкціонованого доступу через механізм захисту.

Метод. Для вирішення поставленого завдання було застосовано метод імітаційного моделювання з використанням принципів моделювання системи масового обслуговування. Даний метод дозволяють отримати основні характеристики системи безпеки інформації від несанкціонованого доступу з обмеженим обсягом буферної пам'яті.

Результати. Розроблено новий алгоритм, моделі та методологію розробки системи безпеки інформації від несанкціонованого доступу, що розглядається як однофазна багатоканальна системи масового обслуговування з обмеженим обсягом буферної пам'яті. Процес одержання результатів моделі було реалізовано системі моделювання General Purpose Simulation System World і проведено порівняльні оцінки основних характеристик системи безпеки інформації щодо різних законів розподілу вихідних параметрів, тобто. при цьому запити несанкціонованого доступу є найпростішими потоками, а час обслуговування підпорядковується експоненційному, постійному та законам Ерлангового розподілу.

Висновки. Проведені експерименти на основі алгоритму та моделі підтвердили очікувані результати при аналізі характеристик системи безпеки інформації від несанкціонованого доступу як однофазної багатоканальної системи масового обслуговування з обмеженим часом очікування запитів у черзі. Ці результати можуть бути використані для практичної побудови нових або модифікації існуючих системи безпеки інформації в мережах обслуговування об'єктів різного призначення. Дана робота є одним з підходів до узагальнення розглянутих проблем для систем з обмеженим обсягом буферної пам'яті. Перспективи подальших досліджень включають в себе дослідження та розробку принципів апаратно-програмної реалізації системи безпеки інформації в мережах обслуговування.

КЛЮЧОВІ СЛОВА: несанкціонований доступ, системи безпеки інформації, інформаційна безпека, системи масового обслуговування, механізм захисту, імітаційна моделювання.

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